

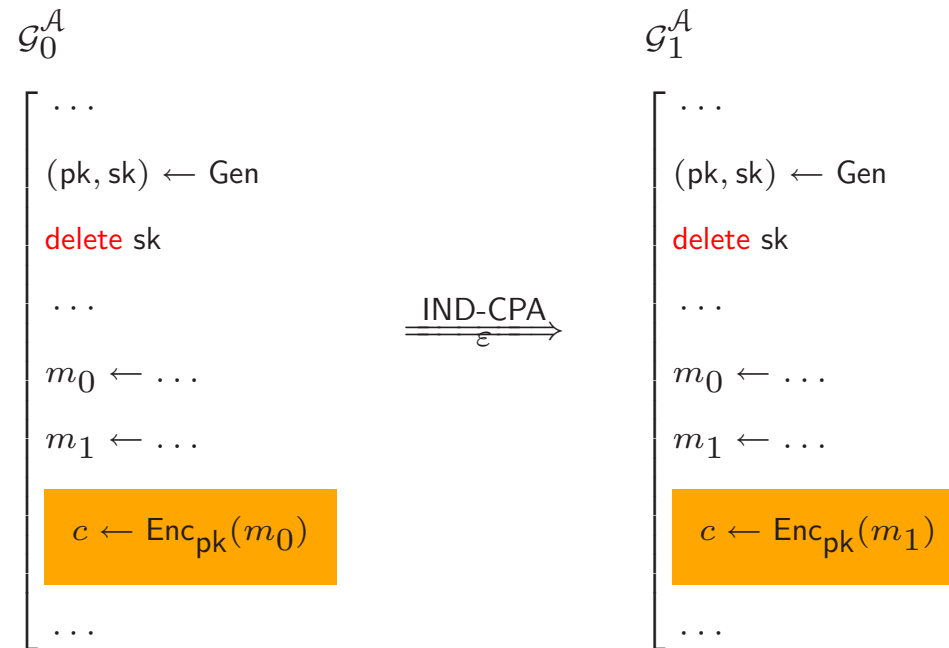
Rewriting Rules for Primitives

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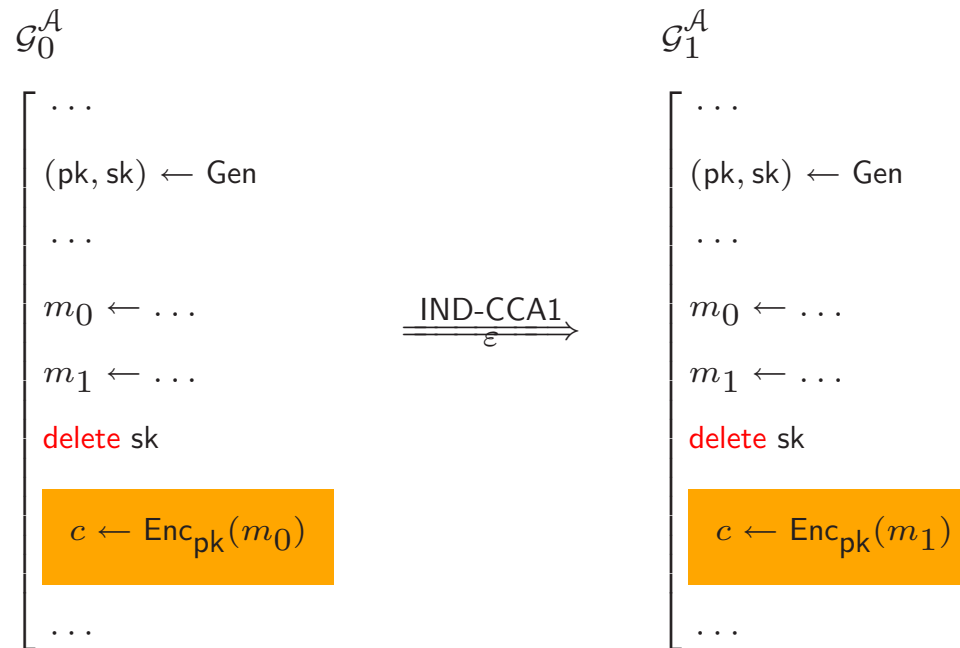
Encryption

IND-CPA security



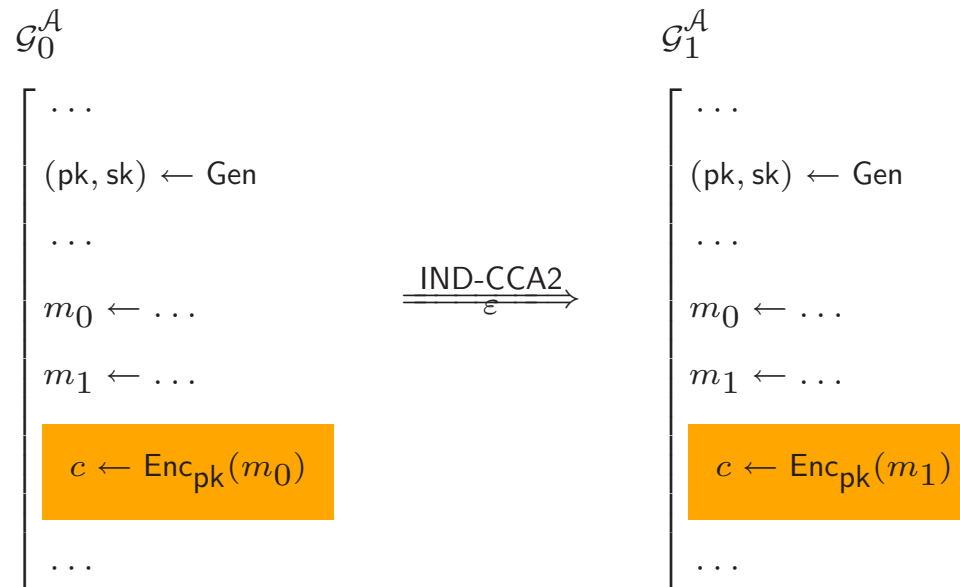
The complexity of the games must be less than t for (t, ε) -IND-CPA secure cryptosystem. In the formal reduction, the whole game is converted to the IND-CPA distinguisher.

IND-CCA1 security



The complexity of the games must be less than t for (t, ε) -IND-CCA1 secure cryptosystem. In the formal reduction, the whole game is converted to the IND-CCA1 distinguisher.

IND-CCA2 security



The reduction is valid if the challenge encryption c is **never** decrypted in the games \mathcal{G}_0 and \mathcal{G}_1 . The complexity of the games must be less than t for (t, ε) -IND-CCA2 secure cryptosystem. In the formal reduction, the whole game is converted to the IND-CCA2 distinguisher.

Message Authentication

MAC security

\mathcal{G}_0^A

```

...
sk ← Gen
...
t1 ← Macsk(m1)
...
tq ← Macsk(mq)
...
if Versk( $\bar{m}$ ,  $\bar{t}$ ) = 0 then return 0
if  $\bar{m} \notin \{m_1, \dots, m_q\}$  then Do something bad
...
    
```

$\xrightarrow[\varepsilon]{\text{MAC}}$

\mathcal{G}_1^A

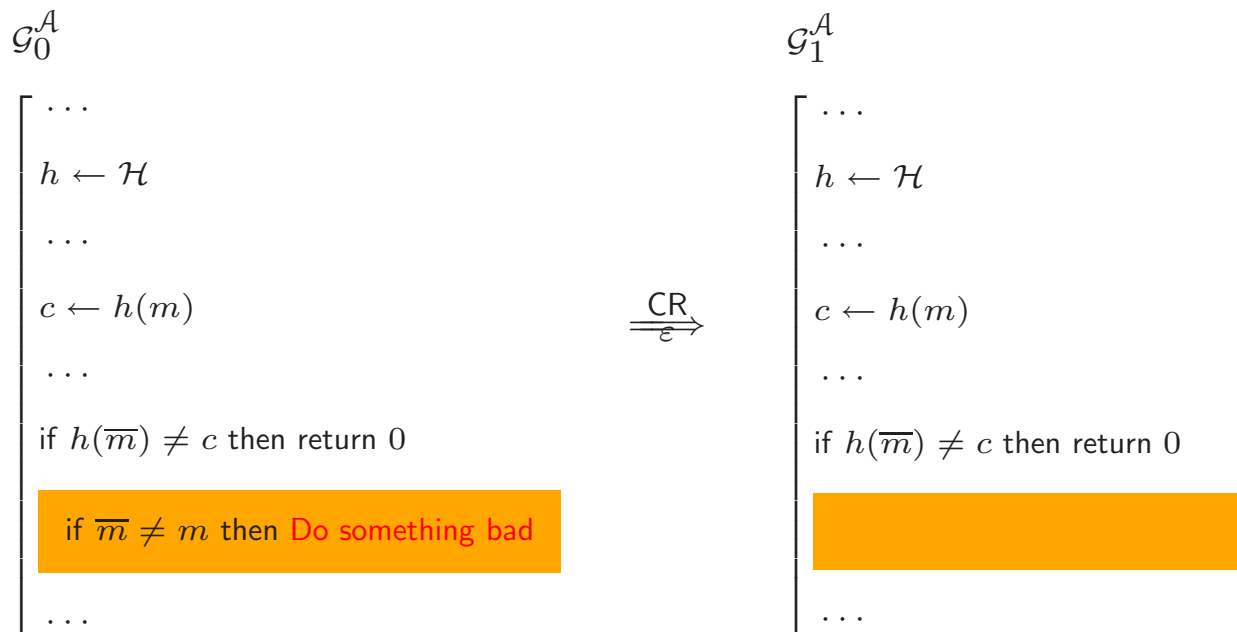
```

...
sk ← Gen
...
t1 ← Macsk(m1)
...
tq ← Macsk(mq)
...
if Versk( $\bar{m}$ ,  $\bar{t}$ ) = 0 then return 0
...
...
    
```

The reduction is valid if sk is used only for computing the $Ver_{sk}(\cdot)$ predicate and for computing $Mac_{sk}(m_1), \dots, Mac_{sk}(m_q)$. The complexity of the games must be less than t for (t, q, ε) -secure message authentication code.

Hash Functions

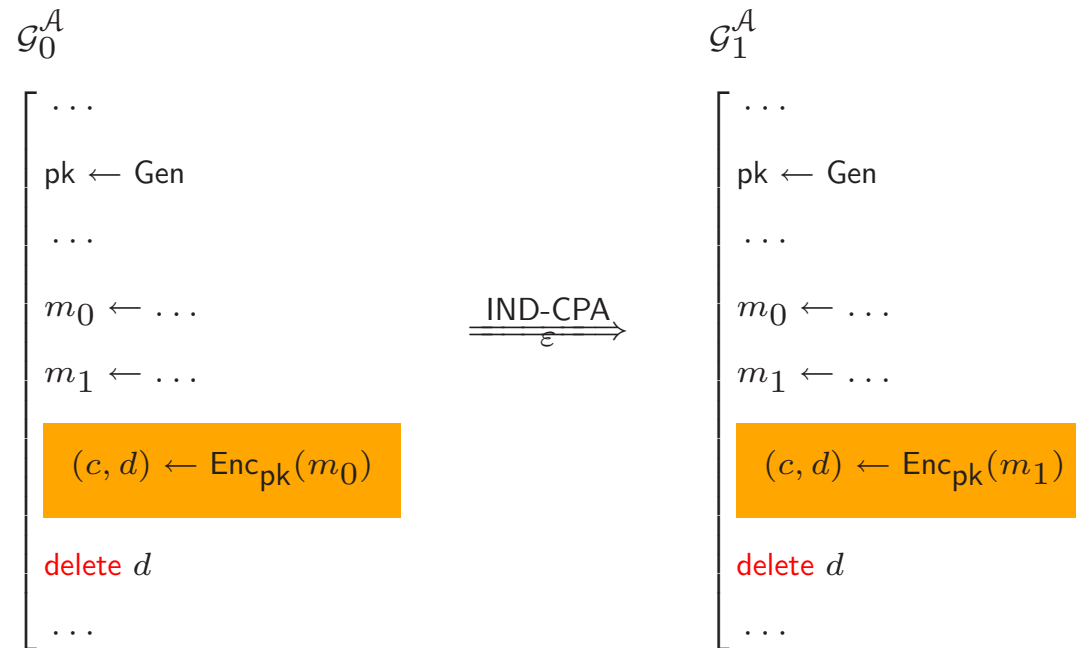
Collision resistance



The complexity of the games must be less than t for (t, ε) -collision resistant hash function family \mathcal{H} .

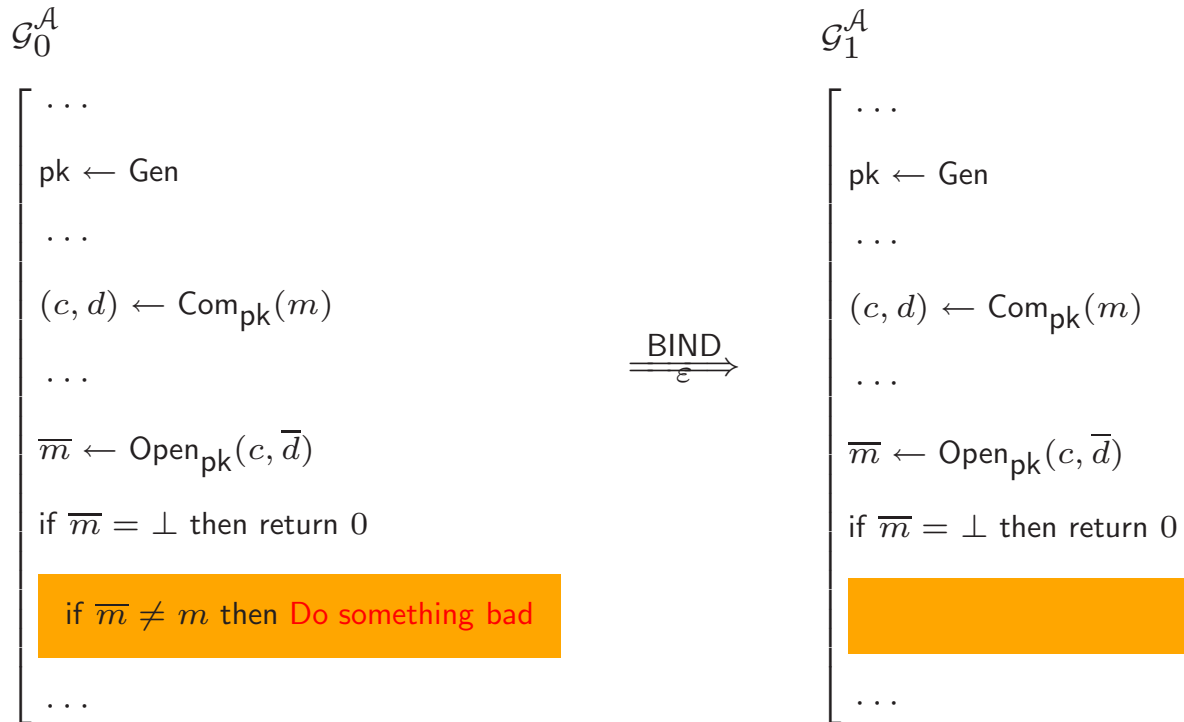
Commitment Schemes

Hiding



The complexity of the games must be less than t for (t, ε) -hiding secure commitment scheme.

Binding



The complexity of the games must be less than t for a (t, ε) -binding commitment scheme.